

# Lower Bound for the Size of Maximal Nontraceable Graphs \*

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## Abstract

Let  $g(n)$  denote the minimum number of edges of a maximal nontraceable graph of order  $n$ . Dudek, Katona and Wojda (2003) showed that  $g(n) \geq \lceil \frac{3n-2}{2} \rceil - 2$  for  $n \geq 20$  and  $g(n) \leq \lceil \frac{3n-2}{2} \rceil$  for  $n \geq 54$  as well as for  $n \in I = \{22, 23, 30, 31, 38, 39, 40, 41, 42, 43, 46, 47, 48, 49, 50, 51\}$ . We show that  $g(n) = \lceil \frac{3n-2}{2} \rceil$  for  $n \geq 54$  as well as for  $n \in I \cup \{12, 13\}$  and we determine  $g(n)$  for  $n \leq 9$ .

**Keywords:** maximal nontraceable, hamiltonian path, traceable, nontraceable, nonhamiltonian

**2000 Mathematics Subject Classification:** 05C38

## 1 Introduction

We consider only simple, finite graphs  $G$  and denote the vertex set, the edge set, the order and the size of  $G$  by  $V(G)$ ,  $E(G)$ ,  $v(G)$  and  $e(G)$ , respectively. The *open neighbourhood* of a vertex  $v$  in  $G$  is the set  $N_G(v) = \{x \in V(G) : vx \in E(G)\}$ . If  $U$  is a nonempty subset of  $V(G)$  then  $\langle U \rangle$  denotes the subgraph of  $G$  induced by  $U$ .

A graph  $G$  is *hamiltonian* if it has a *hamiltonian cycle* (a cycle containing all the vertices of  $G$ ), and *traceable* if it has a *hamiltonian path* (a path containing all the vertices of  $G$ ). A graph  $G$  is *maximal nonhamiltonian* (MNH) if  $G$  is not hamiltonian, but  $G + e$  is hamiltonian for each  $e \in E(\overline{G})$ , where  $\overline{G}$  denotes the complement of  $G$ . A graph  $G$  is *maximal nontraceable* (MNT) if  $G$  is not traceable, but  $G + e$  is traceable for each  $e \in E(\overline{G})$ .

In 1978 Bollobás [1] posed the problem of finding the least number of edges,  $f(n)$ , in a MNH graph of order  $n$ . Bondy [2] had already shown that a MNH

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graph with order  $n \geq 7$  that contained  $m$  vertices of degree 2 had at least  $(3n + m)/2$  edges, and hence  $f(n) \geq \lceil 3n/2 \rceil$  for  $n \geq 7$ . Combined results of Clark, Entringer and Shapiro [3], [4] and Lin, Jiang, Zhang and Yang [7] show that  $f(n) = \lceil 3n/2 \rceil$  for  $n \geq 19$  and for  $n = 6, 10, 11, 12, 13, 17$ . The values of  $f(n)$  for the remaining values of  $n$  are also given in [7].

Let  $g(n)$  denote the minimum number of edges in a MNT graph of order  $n$ . Dudek, Katona and Wojda [5] proved that

$$g(n) \geq \lceil \frac{3n-2}{2} \rceil - 2 \text{ for } n \geq 20$$

and showed, by construction, that

$$g(n) \leq \lceil \frac{3n-2}{2} \rceil \text{ for } n \geq 54$$

as well as for  $n \in I = \{22, 23, 30, 31, 38, 39, 40, 41, 42, 43, 46, 47, 48, 49, 50, 51\}$ .

We prove, using a method different from that in [5], that

$$g(n) \geq \lceil \frac{3n-2}{2} \rceil \text{ for } n \geq 10.$$

We also construct graphs of order  $n = 12, 13$  with  $\lceil \frac{3n-2}{2} \rceil$  edges and thus show that

$$g(n) = \lceil \frac{3n-2}{2} \rceil \text{ for } n \geq 54 \text{ as well as for } n \in I \cup \{12, 13\}.$$

We also determine  $g(n)$  for  $n \leq 9$ .

## 2 Auxilliary Results

In this section we present some results concerning MNT graphs, which we shall use, in the next section, to prove that a MNT graph of order  $n \geq 10$  has at least  $\frac{3n-2}{2}$  edges. The first one concerns the lower bound for the number of edges of MNH graphs. It is the combination of results proved in [2] and [7].

**Theorem 2.1** (*Bondy and Lin, Jiang, Zhang and Yang*) *If  $G$  is a MNH graph of order  $n$ , then  $e(G) \geq \frac{3n}{2}$  for  $n \geq 6$ .*

The following lemma, which we proved in [6], will be used frequently.

**Lemma 2.2** *Let  $Q$  be a path in a MNT graph  $G$ . If  $\langle V(Q) \rangle$  is not complete, then some internal vertex of  $Q$  has a neighbour in  $G - V(Q)$ .*

**Proof.** Let  $u$  and  $v$  be two nonadjacent vertices of  $Q$ . Then  $G + uv$  has a hamiltonian path  $P$ . Let  $x$  and  $y$  be the two endvertices of  $Q$  and suppose no internal vertex of  $Q$  has a neighbour in  $G - V(Q)$ . Then  $P$  has a subpath  $R$  in  $\langle V(Q) \rangle + uv$  and  $R$  has either one or both endvertices in  $\{x, y\}$ . If  $R$  has only one endvertex in  $\{x, y\}$ , then  $P$  has an endvertex in  $Q$ . In either case the path obtained from  $P$  by replacing  $R$  with  $Q$  is a hamiltonian path of  $G$ . ■

The following lemma is easy to prove.

**Lemma 2.3** Suppose  $T$  is a cutset of a connected graph  $G$  and  $A_1, \dots, A_k$  are components of  $G - T$ .

- (a) If  $k \geq |T| + 2$ , then  $G$  is nontraceable.
- (b) If  $G$  is MNT then  $k \leq |T| + 2$ .
- (c) If  $G$  is MNT and  $k = |T| + 2$ , then  $\langle T \cup A_i \rangle$  is complete for  $i = 1, 2, \dots, k$ .

**Proof.** (a) and (b) are obvious. If (c) is not true, then there is an  $i$  such that  $\langle T \cup A_i \rangle$  has two nonadjacent vertices  $x$  and  $y$ . But then  $T$  is a cutset of the graph  $G + xy$  and  $(G + xy) - T$  has  $|T| + 2$  components and hence  $G + xy$  is nontraceable, by (a). ■

The proof of the following lemma is similar to the previous one.

**Lemma 2.4** Suppose  $B$  is a block of a connected graph  $G$ .

- (a) If  $B$  has more than two cut-vertices, then  $G$  is nontraceable.
- (b) If  $G$  is MNT, then  $B$  has at most three cut-vertices.
- (c) If  $G$  is MNT and  $B$  has exactly three cut-vertices, then  $B$  consists of exactly four blocks, each of which is complete.

In [6] we proved some results concerning the degrees of the neighbours of the vertices of degree 2 in a 2-connected MNT graph, which enabled us to show that the average degree of the vertices in a 2-connected MNT graph is at least 3. We now restate those results in a form that is applicable also to MNT graphs which are not 2-connected. (Note that in a 2-connected graph no two vertices of degree 2 are adjacent to one another.)

**Lemma 2.5** If  $G$  is a connected MNT graph and  $v \in V(G)$  with  $d(v) = 2$ , then the neighbours of  $v$  are adjacent. Also, one of the neighbours has degree at least 4 and the other neighbour has degree 2 or at least 4.

**Proof.** Let  $N_G(v) = \{x_1, x_2\}$  and let  $Q$  be the path  $x_1vx_2$ . Since  $N_G(v) \subseteq Q$ , it follows from Lemma 2.2 that  $\langle V(Q) \rangle$  is a complete graph; hence  $x_1$  and  $x_2$  are adjacent.

Since  $G$  is connected and nontraceable, at least one of  $x_1$  and  $x_2$  has degree bigger than 2. Suppose  $d(x_1) > 2$  and let  $z \in N(x_1) - \{v, x_2\}$ . If  $Q$  is the path  $zx_1vx_2$  then, since  $d(v) = 2$ , the graph  $\langle V(Q) \rangle$  is not complete and hence it follows from Lemma 2.2 that  $d(x_1) \geq 4$ . Similarly if  $d(x_2) > 2$ , then  $d(x_2) \geq 4$ . ■

**Lemma 2.6** Suppose  $G$  is a connected MNT graph with distinct nonadjacent vertices  $v_1$  and  $v_2$  such that  $d(v_1) = d(v_2) = 2$ .

- (a) If  $v_1$  and  $v_2$  have exactly one common neighbour  $x$ , then  $d(x) \geq 5$ .
- (b) If  $v_1$  and  $v_2$  have the same two neighbours  $x_1$  and  $x_2$ , then  $N_G(x_1) - \{x_2\} = N_G(x_2) - \{x_1\}$  and  $d(x_1) = d(x_2) \geq 5$ .

**Proof.** (a) Let  $N(v_i) = \{x, y_i\}$ ;  $i = 1, 2$ . It follows from Lemma 2.5 that  $x$  is adjacent to  $y_i$ ;  $i = 1, 2$ . Let  $Q$  be the path  $y_1v_1xv_2y_2$ . Since  $\langle V(Q) \rangle$  is not complete, it follows from Lemma 2.2 that  $x$  has a neighbour in  $G - V(Q)$ . Hence  $d(x) \geq 5$ .

(b) From Lemma 2.5 it follows that  $x_1$  and  $x_2$  are adjacent. Let  $Q$  be the path  $x_2v_1x_1v_2$ .  $\langle V(Q) \rangle$  is not complete since  $v_1$  and  $v_2$  are nonadjacent. Thus it follows from Lemma 2.2 that  $x_1$  has a neighbour in  $G - V(Q)$ . Now suppose  $p \in N_{G-V(Q)}(x_1)$  and  $p \notin N_G(x_2)$ . Then a hamiltonian path  $P$  in  $G + px_2$  contains a subpath of either of the forms given in the first column of Table 1. Note that  $i, j \in \{1, 2\}$ ;  $i \neq j$  and that  $L$  represents a subpath of  $P$  in  $G - \{x_1, x_2, v_1, v_2, p\}$ . If each of the subpaths is replaced by the corresponding subpath in the second column of the table we obtain a hamiltonian path  $P'$  in  $G$ , which leads to a contradiction.

Subpath of $P$	Replace with
$v_ix_1v_jx_2p$	$v_ix_2v_jx_1p$
$v_ix_1Lpx_2v_j$	$v_ix_2v_jx_1Lp$

Table 1

Hence  $p \in N_G(x_2)$ . Thus  $N_G(x_1) - \{x_2\} \subseteq N_G(x_2) - \{x_1\}$ . Similarly  $N_G(x_2) - \{x_1\} \subseteq N_G(x_1) - \{x_2\}$ . Thus  $N_G(x_1) - \{x_2\} = N_G(x_2) - \{x_1\}$  and hence  $d(x_1) = d(x_2)$ . Now let  $Q$  be the path  $px_1v_1x_2v_2$ . Since  $\langle V(Q) \rangle$  is not complete, it follows from Lemma 2.2 that  $x_1$  or  $x_2$  has a neighbour in  $G - V(Q)$ . Hence  $d(x_1) = d(x_2) \geq 5$ . ■

**Lemma 2.7** *Suppose  $G$  is a connected MNT graph of order  $n \geq 6$  and that  $v_1, v_2$  and  $v_3$  are vertices of degree 2 in  $G$  having the same neighbours,  $x_1$  and  $x_2$ . Then  $G - \{v_1, v_2, v_3\}$  is complete and hence  $e(G) = \frac{1}{2}(n^2 - 7n + 24)$ .*

**Proof.** The set  $\{x_1, x_2\}$  is a cutset of  $G$ . Thus according to Lemma 2.3  $G - \{v_1, v_2, v_3\} = K_{n-3}$ . Hence  $e(G) = \frac{1}{2}(n-3)(n-4) + 6$ . ■

By combining the previous three results we obtain

**Theorem 2.8** *Suppose  $G$  is a connected MNT graph without vertices of degree 1 or adjacent vertices of degree 2. If  $G$  has order  $n \geq 7$  and  $m$  vertices of degree 2, then  $e(G) \geq \frac{1}{2}(3n + m)$ .*

**Proof.** If  $G$  has three vertices of degree 2 having the same two neighbours then, by Lemma 2.7,  $m = 3$  and

$$e(G) = \frac{1}{2}(n^2 - 7n + 24) \geq \frac{1}{2}(3n + m) \text{ when } n \geq 7.$$

We now assume that  $G$  does not have three vertices of degree 2 that have the same two neighbours. Let  $v_1, \dots, v_m$  be the vertices of degree 2 in  $G$  and let  $H = G - \{v_1, \dots, v_m\}$ . Then by Lemmas 2.5 and 2.6 the minimum degree,  $\delta(H)$  of  $H$  is at least 3. Hence

$$e(G) = e(H) + 2m \geq \frac{3}{2}(n - m) + 2m = \frac{1}{2}(3n + m).$$

■

### 3 The minimum size of a MNT graph

Our aim is to determine the exact value of  $g(n)$ . By consulting the Atlas of Graphs [8], one can see, by inspection, that  $g(2) = 0$ ,  $g(3) = 1$ ,  $g(4) = 2$ ,  $g(5) = 4$ ,  $g(6) = 6$  and  $g(7) = 8$  (see Fig. 3).

We now give a lower bound for  $g(n)$  for  $n \geq 8$ .

**Theorem 3.1** *If  $G$  is a MNT graph of order  $n$ , then*

$$e(G) \geq \begin{cases} 10 & \text{if } n = 8 \\ 12 & \text{if } n = 9 \\ \frac{3n-2}{2} & \text{if } n \geq 10. \end{cases}$$

**Proof.**

If  $G$  is not connected, then  $G = K_k \cup K_{n-k}$ , for some positive integer  $k < n$  and then, clearly,  $e(G) > \frac{3n-2}{2}$  for  $n \geq 8$ . Thus we assume that  $G$  is connected.

We need to prove that the sum of the degrees of the vertices of  $G$  is at least  $3n - 2$ . In view of Theorem 2.8, we let

$$M = \{v \in V(G) \mid d(v) = 2 \text{ and no neighbour of } v \text{ has degree } 2\}.$$

The remaining vertices of degree 2 can be dealt with simultaneously with the vertices of degree 1. We let

$$S = \{v \in V(G) - M \mid d(v) = 2 \text{ or } d(v) = 1\}.$$

If  $S = \emptyset$ , then it follows from Theorem 2.8 that  $e(G) \geq \frac{1}{2}(3n + m)$ . Thus we assume that  $S \neq \emptyset$ .

We observe that, if  $H$  is a component of the graph of  $\langle S \rangle$ , then either  $H \cong K_1$  or  $H \cong K_2$  and  $N_G(H) - V(H)$  consists of a single vertex, which is a cut-vertex of  $G$ .

An example of such a graph  $G$  is depicted in the figure below.

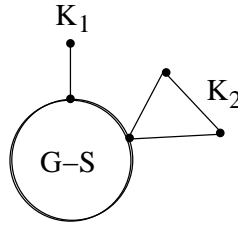


Fig. 1

Let  $s = |S|$ . By Lemma 2.4 the graph  $\langle S \rangle$  has at most three components. We thus have three cases:

**CASE 1.**  $\langle S \rangle$  has exactly three components, say  $H_1, H_2, H_3$ :

In this case the neighbourhoods of  $H_1, H_2, H_3$  are pairwise disjoint; hence  $G$  has three cut-vertices. Hence it follows from Lemma 2.4 that  $G - S$  is a complete graph of order at least 3. Furthermore, for every possible value of  $s$ , the number of edges in  $G$  incident with the vertices in  $S$  is  $2s - 3$ . Thus

$$e(G) = \binom{n-s}{2} + 2s - 3 \text{ for } s = 3, 4, 5 \text{ or } 6; s \leq n - 3.$$

An easy calculation shows that, for each possible value of  $s$ ,

$$e(G) \geq \begin{cases} 10 & \text{if } n = 8 \\ 12 & \text{if } n = 9 \\ \frac{3n-2}{2} & \text{if } n \geq 10. \end{cases}$$

This case is a Zelinka Type II construction, cf. [9]. The graphs of smallest size of order 8 and 9 given by this construction are depicted in Fig. 3.

**CASE 2.**  $\langle S \rangle$  has exactly two components, say  $H_1, H_2$ :

In this case the number of edges in  $G$  incident with the vertices in  $S$  is  $2s - 2$ .

**Subcase 2.1.**  $N_G(H_1) = N_G(H_2)$ :

Then it follows from Lemma 2.3 that  $G - S$  is a complete graph. Hence

$$e(G) = \binom{n-s}{2} + 2s - 2 \text{ for } s = 2, 3 \text{ or } 4.$$

Thus

$$e(G) \geq \begin{cases} 12 & \text{if } n = 8 \\ 16 & \text{if } n = 9 \\ \frac{3n-2}{2} & \text{if } n \geq 10 \end{cases}$$

This case is a Zelinka Type I construction, cf. [9].

**Subcase 2.2.**  $N_G(H_1) \neq N_G(H_2)$ :

Let  $N_G(H_i) = y_i$ ,  $i = 1, 2$  and  $y_1 \neq y_2$ .

If  $y_1 y_2 \notin E(G)$  then  $G + y_1 y_2$  has a hamiltonian path  $P$ . But then  $P$  has one endvertex in  $H_1$  and the other in  $H_2$  and contains the edge  $y_1 y_2$ ; hence  $V(G - S) = \{y_1, y_2\}$ . But then  $G$  is disconnected. This contradiction shows that  $y_1 y_2 \in E(G)$ .

Now  $G - S$  is not complete, otherwise  $G$  would be traceable. Since  $G + vw$ , where  $v$  and  $w$  are nonadjacent vertices in  $V(G - S)$ , contains a hamiltonian path with one endvertex in  $H_1$  and the other in  $H_2$  and  $y_1 y_2 \in E(G)$ , it follows that  $(G - S) + vw$  has a hamiltonian cycle. Hence  $G - S$  is either hamiltonian or MNH. We consider these two cases separately:

**Subcase 2.2.1.**  $G - S$  is hamiltonian:

Then no hamiltonian cycle in  $G - S$  contains  $y_1y_2$ , otherwise  $G$  would be traceable. Thus  $d_{G-S}(y_i) \geq 3$  for  $i = 1, 2$ .

It also follows from Lemma 2.3 that no vertex  $v \in M$  can be adjacent to both  $y_1$  and  $y_2$  since the graph  $\langle V(H_i) \cup T \rangle$ , where  $T = \{y_1, y_2\}$  is not complete, for  $i = 1, 2$ . If  $v \in M$  is adjacent to one of the  $y_i$ 's for  $i = 1, 2$ , say  $y_1$ , then, since the neighbours of  $v$  are adjacent, it follows that  $d_{G-M-S}(y_1) \geq 3$ .

It follows from our definition of  $M$  and  $S$  that  $N_G(M) \cap S = \emptyset$ . Since  $G - M$  is not a complete graph, it follows from Lemma 2.7 that  $M$  does not have three vertices that have the same neighbourhood in  $G$ . Hence, by Lemmas 2.5 and 2.6, the minimum degree of the graph  $G - M - S$  is at least 3.

Now, for  $n \geq 8$

$$\begin{aligned} e(G) &= e(G - M - S) + 2m + 2s - 2 \\ &\geq \frac{1}{2}(3(n - m - s)) + 2m + 2s - 2 \\ &= \frac{1}{2}(3n + m + s - 4) \\ &\geq \frac{3n - 2}{2}, \text{ since } s \geq 2. \end{aligned}$$

**Subcase 2.2.2.**  $G - S$  is nonhamiltonian:

Then  $G - S$  is MNH (as shown above); hence it follows from Theorem 2.1, that  $e(G - S) \geq \frac{3}{2}(n - s)$  for  $n - s \geq 6$ .

Thus, for  $n - s \geq 6$  and  $n \geq 8$

$$\begin{aligned} e(G) &= e(G - S) + 2s - 2 \\ &\geq \frac{1}{2}(3(n - s)) + 2s - 2 \\ &= \frac{1}{2}(3n + s - 4) \\ &\geq \frac{3n - 2}{2}, \text{ since } s \geq 2. \end{aligned}$$

From [7] we have

$$e(G - S) \geq \begin{cases} 6 & \text{for } n - s = 5 \\ 4 & \text{for } n - s = 4. \end{cases}$$

Thus

$$e(G) \geq \begin{cases} 12 & \text{for } n = 9 \text{ and } n - s = 5 \\ 10 & \text{for } n = 8 \text{ and } n - s = 5 \text{ or } n - s = 4. \end{cases}$$

The smallest MNH graphs  $F_4$  and  $F_5$  of order 4 and 5 respectively, are depicted in Fig. 2; cf. [7]. The graphs  $G_8$  and  $G_9$  (see Fig. 3) are obtained,

respectively, by using  $F_4$  with  $s = 4$  or  $F_5$  with  $s = 3$ , and  $F_5$  with  $s = 4$ .

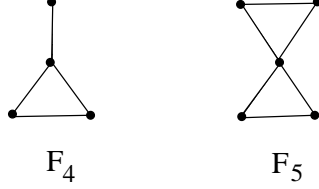


Fig. 2

**CASE 3.**  $\langle S \rangle$  has exactly one component, say  $H$ :

Since

$$\sum_{v \in S} d_G(v) = 3s - 2, \text{ for } s = 1, 2$$

it follows that

$$\begin{aligned} e(G) &= e(G - M) + 2m \\ &= \frac{1}{2} \left( \sum_{v \in V(G-M)-S} d_{G-M}(v) + \sum_{v \in S} d_{G-M}(v) \right) + 2m \\ &\geq \frac{1}{2} (3(n - m - s) + 3s - 2) + 2m \\ &= \frac{1}{2} (3n + m - 2) \\ &\geq \frac{3n - 2}{2}. \end{aligned}$$

■

From the previous theorem we have  $g(8) = 10$ ,  $g(9) = 12$  and  $g(n) \geq \lceil \frac{3n-2}{2} \rceil$  for  $n \geq 10$ . The MNT graphs  $G_n$  of order  $n$  with  $g(n)$  edges, for  $n \leq 9$  are given in Fig. 3.

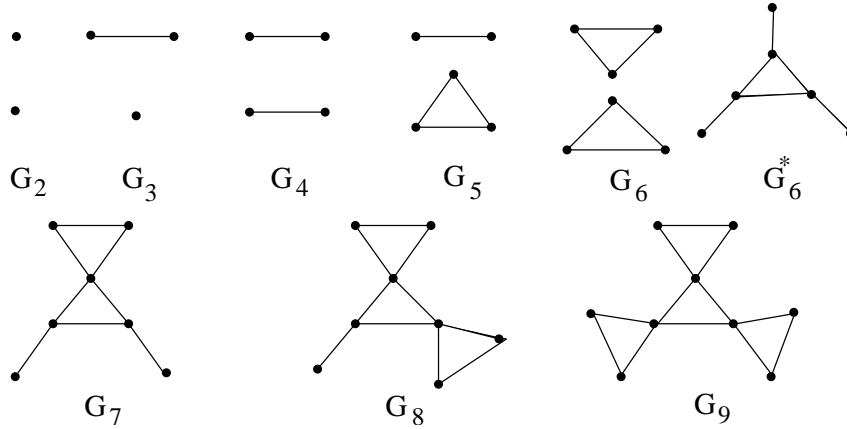


Fig. 3



In [5] Dudek, Katona and Wojda constructed, for every  $n \geq 54$  as well as for every  $n \in I = \{22, 23, 30, 31, 38, 39, 40, 41, 42, 43, 46, 47, 48, 49, 50, 51\}$ , a MNT graph of size  $\lceil \frac{3n-2}{2} \rceil$  in the following way: Consider a cubic MNH graph  $G$  with the property that

- (1) there is an edge  $y_1y_2$  of  $G$ , such that  $N(y_1) \cap N(y_2) = \emptyset$ , and
- (2)  $G + e$  has a hamiltonian cycle containing  $y_1y_2$  for every  $e \in E(\overline{G})$ .

Now take two graphs  $H_1$  and  $H_2$ , with  $H_1 \cong K_1$  and  $H_2 \cong K_1$  or  $H_2 \cong K_2$  and join each vertex of  $H_i$  to  $y_i$ ;  $i = 1, 2$ . The new graph is a MNT graph of order  $v(G) + 2$  and size  $e(G) + 2$  or of order  $v(G) + 3$  and size  $e(G) + 4$ .

It follows from results in [3] and [4] that for every even  $n \geq 52$  as well as for  $n \in \{20, 28, 36, 38, 40, 44, 46, 48\}$  there exists a cubic MNH graph of order  $n$  that satisfies (1) and (2). Thus this construction provides MNT graphs of order  $n$  and size  $\lceil \frac{3n-2}{2} \rceil$  for every  $n \geq 54$  as well as for every  $n \in I$ .

We determined, by using the Graph Manipulation Package developed by Siqinfu and Sheng Bau\*, that the Petersen graph also satisfies the above property. Hence, according to the above construction, there are also MNT graphs of order  $n$  and size  $\lceil \frac{3n-2}{2} \rceil$  for  $n = 12, 13$ .

Thus  $g(n) = \lceil \frac{3n-2}{2} \rceil$  for  $n \geq 54$  as well as for every  $n \in I \cup \{12, 13\}$ .

It remains an open problem to find  $g(n)$  for  $n = 10, 11$  and those values of  $n$  between 13 and 54 which are not in  $I$ .

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